







Private-Safe (Logic-based) Decision Systems for Energy Assignment in Agricultural Cooperatives

<u>Luciana Fidilio-Allende</u> Joaquín Arias

Grupo de Inteligencia Artificial de la URJC Center for Intelligent Information Technologies (CETINIA) Móstoles, Madrid

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- Automated Decision Makers, even if they are value-aware, can only be trustworthy if they are capable of explaining their decisions (XAI).
- Logic-based systems, such as s(CASP) [3], a goal-directed execution of Answer Set Programming (ASP) [7], can provide justifications.
 - But the justifications (and the ASP models themselves) may expose sensitive information (violating privacy and/or legislation).



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```
Example from [9]
```

```
Justification for ?- person(mary)
```

```
person(mary) :-
person(X) :- ustaff(X).
ustaff(X) :- professor(X).
                               ustaff(mary) :-
professor(mary).
                                   professor(mary).
```



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Example from [9]

```
person(X) :- ustaff(X). 1
ustaff(X) :- professor(X). 2
professor(mary). 3
```

Justification for ?- person(mary)

```
'person' holds (for mary), because
'ustaff' holds (for mary), because
'professor' holds (for mary).
```



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Limitations of forgetting proposals/operators

- (i) They have technical limitations.
- (ii) They are focused on propositional logic, or need to ground the variables.



Limitations of state-of-the-art Forgetting techniques

	(UP)	(SP)	Loops	Commutative	Variables	Constraints
f_{SU} [8]	Yes	No	Yes	No	No	No
f_{SP} [6]	Yes	Limited	No	No	No	No
f_{SP}^{*} [4]	Yes	Limited	Yes	Yes	No	No
f_{AC} [5]	Yes	Yes	Yes	Yes	No	No

- (UP): Uniform Persistence means that the original program and the forgetting result are equivalent even if we add new facts.
- (SP): Strong Persistence is similar to (UP) but adding new rules.
- Loops: Deal with even/odd loops (by adding auxiliary predicates).
- Commutative: Allow iterative application, regardless of the order.

Limitations of state-of-the-art Forgetting techniques (cont.)

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f_{SP} [6]	Yes	Limited	No	No	No	No
f_{SP}^{*} [4]	Yes	Limited	Yes	Yes	No	No
f_{AC} [5]	Yes	Yes	Yes	Yes	No	No
f_{CASP}	Yes	Yes	Yes	Yes	WiP	WiP

Our proposal f_{CASP}

based on f_{SU}

- It is simple, iterable, and commutable.
- Uses the compiler of s(CASP) to generate dual rules.
- Supports loops and denials.
 - ...and we believe that it can satisfy SP and, in the future, support Variables and Constraints.



Background: ASP

 Answer Set Programming (ASP) is based on the stable model semantics [7], supporting non-stratified negation:

- In this work, we extended ASP with **double default negations** [10]: The clause p:- not not p. generates two models: {p} and {}.
 - Double negations can be modeled as even loops. For example, the predicate p :- not not p is transformed into:

```
p :- not neg_p.
neg_p :- not p.
```



Background: s(CASP)

- It is a top-down, goal-directed interpreter of ASP with Constraints [3].
- Can provide justifications (in natural language) [1].
 - They can be manipulated (to hide sensitive information) using the directive #show and the flag --short.
- It solves negated atoms (not p) against the dual rules of the program (the negation of the rules present in the program) [2]. E.g.:

```
% Original program
  p(0).
p(X) := q(X), \text{ not } t(X,Y).
```

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```
1 % Original program
                                              % Dual rules
_{2} p(0).
                                              not p(X) := not p1(X), not p2(X).
p(X) := q(X), \text{ not } t(X,Y).
                                             not p1(X) := X = 0.
                                             not p2(X) := forall(Y, not <math>p2_{(X,Y)}).
                                             not p2_{(X,Y)} := not q(X).
                                            6 not p2(X,Y) := q(X), t(X,Y).
```



Background: The forgetting technique f_{CASP}

 f_{CASP} consists on five steps that can be iteratively repeated to forget multiple predicates (consider we want to forget the predicate p):

- Add auxiliary predicates due to even loops, facts and/or missing predicate.
 If p is part of an even loop, not p is replaced, and neg_x :- not p is added.
- 2. Generate the simplified dual rule(s) using s(CASP).

```
% Clauses of p
p :- t, not u.
p :- not r.

4 not p_1 :- u.
5 not p_2 :- r.
1 % s(CASP) dual rules
1 % Simplified dual rule(s)
2 not p :- not p_2.
2 not p :- not t, r.
3 not p_1 :- u.
5 not p_2 :- r.
```

- 3. Forget the predicate and its negation.
- 4. Clean true/false and add double negations to preserve even loops.

Repeat steps 1 to 4 to forget the next predicate...

5. (Optional) Transform double negations into even loops.



Energy assignment in agricultural cooperatives





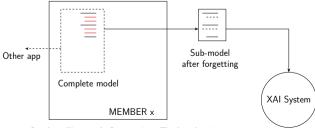
Ethical values in energy management

- In agricultural cooperatives, farmers share their assets to obtain a better and cheaper access to resources. The sharing of locally-generated energy can be a stable alternative supply in rural areas.
- Energy assignment can encourage better practices if we base the distribution criteria of the generated energy on values.



- We can consider values based on the 10 CAP objectives as criteria.
- For example, we can consider how fairly the farmer's workers are paid.

- We propose an automated decision-making system for energy assignment in agricultural cooperatives using fair income as criteria.
 - Its decisions must be fair.
 - To trust a decision, a justification is required (XAI).
 - Additionally, the members of the cooperative may want to preserve business secrets (e.g., salary complements).



- We propose an automated decision-making system for energy assignment in agricultural cooperatives using fair income as criteria.
 - Its decisions must be fair.
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 - Additionally, the members of the cooperative may want to preserve business secrets (e.g., salary complements).

```
% Original clauses
                                            % Result after forgetting
salary(eric, Salary):-
                                         salary(eric, Salary):-
 base salarv(eric, SO),
                                              S0 = 1200.
 distance_home_work(eric, S1),
                                         S1 = 100.
 has_children(eric, S2),
                                         5 	 S2 = 100.
 Salarv is S0 + S1 + S2.
                                              Salarv is S0 + S1 + S2.
```



• In other scenarios the clauses involve even loops.

```
% Original clauses
   over 40 bea :-
     not neg_over_40_bea.
   neg_over_40_bea:-
     not over_40_bea.
6
   generational_renewal(bea, 0):-
      over 40 bea.
   generational_renewal(bea, 100):-
10
     not over 40 bea.
11
   salary(bea, Salary):-
     base_salary(bea, S0),
13
      generational_renewal(bea, S1),
14
     holiday_worked(bea, S2),
15
     Salarv is S0 + S1 + S2.
16
```



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```
% Original clauses
   over 40 bea :-
     not neg over 40 bea.
   neg_over_40_bea:-
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10
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11
   salary(bea, Salary):-
12
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13
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14
     holiday_worked(bea, S2),
15
     Salary is S0 + S1 + S2.
16
```



• In other scenarios the clauses involve even loops.

```
% Original clauses
                                                   % Result after forgetting
   over 40 bea :-
                                                   neg_1 :- not neg_2.
     not neg over 40 bea.
                                                   neg 2 :- not neg 1.
   neg_over_40_bea:-
                                                4
     not over_40_bea.
                                                   salary(bea, Salary):-
                                                     SO = 900.
6
   generational_renewal(bea, 0):-
                                                     neg_2,
      over 40 bea.
                                                     S1 = 0,
   generational_renewal(bea, 100):-
                                                     S2 = 0.
     not over_40_bea.
                                               10
                                                     Salarv is S0 + S1 + S2.
10
                                                   salarv(bea, Salarv):-
                                               11
11
   salary(bea, Salary):-
                                                     S0 = 900.
                                               12
     base_salary(bea, S0),
                                                     neg_1,
13
                                               13
     generational_renewal(bea, S1),
                                                     S1 = 100.
14
                                               14
     holiday_worked(bea, S2),
                                                     S2 = 0,
15
                                               15
     Salarv is S0 + S1 + S2.
                                                     Salarv is S0 + S1 + S2.
16
                                               16
```



• Let's see the justifications for ?- salary(bea).

```
% Answer 1
                                                      % Answer 2
salary(bea, 1000) :-
                                                      salary(bea,900) :-
    base_salary(bea,900),
                                                          base_salary(bea,900),
    generational_renewal(bea,100) :-
                                                          generational_renewal(bea,0) :-
        not over 40 bea :-
                                                              over 40 bea :-
            neg_over_40_bea :-
                                                                  not neg_over_40_bea :-
              chs(not over 40 bea).
                                                                     chs(over 40 bea).
    holiday worked(bea.0).
                                                          holiday worked(bea.0).
    1000 is 900+100+0.
                                                          900 \text{ is } 900+0+0
```



• Let's see the justifications for ?- salary(bea).

```
% Answer 1
                                                     % Answer 1
salary(bea, 1000) :-
                                                     salary(bea, 1000) :-
    base_salary(bea,900),
                                                       neg_1 :-
    generational_renewal(bea,100) :-
                                                           not neg_2 :-
        not over 40 bea :-
                                                               chs(neg 1).
                                                       1000 is 900+100+0.
            neg_over_40_bea :-
              chs(not over 40 bea).
    holiday worked(bea.0).
    1000 is 900+100+0.
```



Conclusions

- We have presented a (simple) use case of the forgetting operator f_{CASP}, modelling a decision system for a fair distribution of the energy in an agricultural cooperative.
- In this use case, the application of the operator has made possible to preserve the confidentiality of the farmers
 - ... while maintaining the explainability of the model.



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Future Work

- Expand f_{CASP} to support generic CASP programs.
- Test the technical limits of the operator.
- Provide a formal proof of the f_{CASP} algorithm's correctness.



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